

Non-cooperative Support for the Asymmetric Nash Bargaining Solution

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Abstract

Our work contributes to the game-theoretic analysis of bargaining by providing additional non-cooperative support to the well-known Nash bargaining solution. In particular, in the present paper we study a model of non-cooperative multilateral bargaining with a very general proposer selection protocol and set of feasible payoffs. In each period of the bargaining game, one out of n players is recognized as the proposer according to an irreducible Markov process. The proposer offers a particular element of the convex set of feasible payoffs. If all players accept the offer, it is implemented. If a player rejects the offer, with some probability the negotiations break down and with the remaining probability the next period starts. We show that subgame perfect equilibria in stationary strategies exist and we fully characterize the set of such equilibria. Our main result is that in the limit, as the exogenous risk of breakdown goes to zero, stationary subgame perfect equilibrium payoffs converge to the weighted Nash bargaining solution with the stationary distribution of the Markov proposer selection process as the weight vector. This result generalizes recent findings which are based on time-invariant recognition probabilities or a fixed order of proposers. The proof requires more than mere technical generalizations of those existing results.

KEYWORDS: Nash bargaining solution, subgame perfect equilibrium, stationary strategies, Markov process.

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1 Introduction

This paper contributes to the Nash program of supporting solution concepts from cooperative game theory by obtaining them as equilibrium outcomes of suitably constructed non-cooperative games. More specifically, we will be concerned with the asymmetric Nash bargaining solution. Consider a situation where two players receive a given pair of payoffs if they disagree, but may obtain any element of a convex set of other (superior) payoff pairs if they mutually agree on one such element. The Nash bargaining solution (NBS) is that payoff pair which maximizes the product of players' gains over their disagreement payoff. Nash (1950) showed that this is the unique bargaining solution satisfying the axioms of scale invariance, symmetry, efficiency, and independence of irrelevant alternatives. One can generalize the NBS by assigning different weights to the players. The asymmetric Nash bargaining solution (ANBS) is that payoff pair which maximizes a weighted product of players' gains over their disagreement payoff, see Kalai (1977).

The NBS is used to gain insights on a wide variety of problems in economics. For instance, Bester (1993) compares the effects of different pricing mechanisms on price and quality of a product. In particular, posted pricing is compared to bargaining between a buyer and a seller. In the latter case, the ANBS is taken to be the outcome of the bargaining interaction.

Another common application is wage bargaining between a firm and a union: Firm owners and workers can agree to produce and hence create a surplus. A part of the surplus goes to the workers as their wage, and the rest goes to the shareholders. If, however, the two sides cannot find an agreement, the workers may strike or the firm may shut down so that no surplus is generated. In a seminal paper on wage bargaining, Grout (1984) studies the effect of different legal frameworks on wage bargaining. Throughout the analysis, it is assumed that bargaining between the firm and the union leads to the outcome predicted by the ANBS.

An application to macroeconomic policy making is given by Alesina (1987): In a democracy with two political parties, each party represents voters with different preferences over macroeconomic policy outcomes. Elections take place regularly, but the probability that one or the other party wins is fixed exogenously. The parties may either implement their most preferred policy whenever they are in power or they may agree on a compromise which will be implemented irrespective of who is in power at any specific point in time. Under appropriate assumptions on utility functions, making a compromise generates a surplus. The ANBS is used to predict the division of this surplus; the weights for the ANBS are given by the exogenous probability of winning an election.

The prominent use of the NBS in applications highlights the need for strong non-cooperative underpinnings of this concept. In the case of the ANBS, it is imperative to

examine the non-cooperative or strategic sources of players' "bargaining power" which is borne out in the weight vector of the ANBS. Nash (1953) presents a non-cooperative demand game with two players who are uncertain about which payoff pairs are feasible. In the limit as the uncertainty vanishes, equilibrium payoffs converge to those predicted by the NBS. Carlsson (1991) takes a similar approach, but with a different source of uncertainty: While the set of feasible payoffs is known to both players, their actions are subject to noise. If players make demands which do not exhaust the available surplus, the remainder is distributed according to an exogenously fixed rule. In the limit as the noise vanishes, there is a unique efficient equilibrium. The payoff pair is a particular ANBS; the bargaining weights are determined by the exogenous division rule.

In a seminal paper, Rubinstein (1982) provides a non-cooperative game in which two players negotiate on the division of a pie. The players take turns acting as the proposer. The division of the pie in the unique subgame perfect equilibrium depends upon how strongly players prefer current over future payoffs. In the limit as players become perfectly patient, the equilibrium division converges to the NBS. In their discussion of cooperative and non-cooperative approaches to bargaining, Binmore, Rubinstein and Wolinsky (1986) obtain the NBS in the limit if either players' impatience or the risk of an exogenous breakdown of the negotiations is vanishing.

Although the relationship between cooperative and non-cooperative approaches to bargaining are well understood for the case of two players, such is far less the case when more than two players are involved in the negotiation process. While it is straightforward to generalize the NBS to n players, the extension of its non-cooperative justification has turned out to be a much more difficult problem. Krishna and Serrano (1996) make use of Lensberg's (1988) stability (consistency) property. They design a non-cooperative bargaining protocol in which players can exit after partial agreements. This game has a unique subgame perfect equilibrium and the payoffs implied by that equilibrium converge to the NBS as the discount factor goes to one. Chae and Yang (1994) obtain uniqueness of perfect equilibrium and convergence to the NBS in a game where a proposer negotiates with one responder at a time. In both papers, the results come at the cost of allowing partial agreements, rather than requiring unanimous consent to a comprehensive proposal.

More recently, support results for the ANBS have been given by Miyakawa (2006) and Laruelle and Valenciano (2007). They consider a bargaining game where the proposer in each period is drawn from an invariant probability distribution and unanimous consent is required. The stationary equilibrium payoffs turn out to converge to the ANBS with that probability distribution as the weight vector. Our analysis covers this result as a special case. Another special case is a fixed order of proposers, as analyzed by Kultti and Vartiainen (2007), who also show that differentiability of the payoff set's Pareto frontier is

essential for the convergence result if there are at least three players.

Other strands of the bargaining literature consider much more general bargaining protocols. For instance, in their analysis of uniqueness and efficiency of equilibria in bargaining games, Merlo and Wilson (1995) assume that both the size of the cake to be divided and the order in which players propose and respond follow a Markov process. Kalandrakis (2004) examines no-delay equilibria in stationary strategies under a Markov selection protocol, where agreement does not necessarily require unanimous consent of all players, but only of those within a winning coalition. The set of such equilibria is characterized and shown to be non-empty. Herings and Predtetchinski (2007) study a game with Markov recognition probabilities, where the set of feasible payoffs is one-dimensional. While studying delay or inefficiencies in bargaining games, other authors have used proposer protocols following stochastic processes, see for instance Cho and Duggan (2005) and Hyndman and Ray (2007). In these papers, the stochastic process is not even required to have the Markov property, so the choice of the proposer may depend on aspects of history other than the identity of the previous proposer. Since more general selection protocols are used in much of the bargaining literature, we find it important to extend this approach to the support results for the NBS.

In this paper we take a general approach towards multilateral bargaining. We aim at results for the case with n players, a general set of feasible payoffs, and a general bargaining protocol. An informal description of the game we consider in this paper is as follows.

In the first period of an infinitely repeated bargaining game, the identity of the proposer is completely arbitrary. In each subsequent period, one out of the n players is recognized as the proposer according to an irreducible Markov process. Upon recognition, the proposer offers a particular element of a convex and comprehensive set of feasible payoffs. If all players accept the offer, it is implemented. If a player rejects the offer, with some exogenously given and constant breakdown probability the game ends, whereas with the complementary probability the next period starts.

We show that subgame perfect equilibria in stationary strategies exist and we characterize the set of such equilibria. We then study the limit of an arbitrary sequence of such equilibria corresponding to a sequence of vanishing breakdown probabilities. We show that in the limit all players make the same proposal. Our main result is that in the limit this common proposal coincides with the ANBS with the stationary distribution of the Markov proposer selection process as the weight vector. Hence, equilibrium payoffs depend only on the set of feasible payoffs and the stationary distribution associated with the matrix of transition probabilities.

The proof of our result goes well beyond mere technical generalizations of existing proof strategies. Since the reservation payoff of a responding player depends on the identity

of the current proposer, reservation payoffs cannot be expressed by a single vector, but correspond to a matrix. For any value of the exogenous breakdown probability, we consider the vectors corresponding to the difference between the equilibrium proposals of a pair of players. We show that in the limit as the breakdown probability vanishes, these vectors span an $(n-1)$ -dimensional supporting hyperplane to the set of feasible payoffs at the point corresponding to the common limit proposal of the players. Finally, we demonstrate that the unique normal vector to this supporting hyperplane is proportional to the gradient of the asymmetric Nash product with weights equal to the stationary distribution associated with the matrix of transition probabilities.

The rest of the paper is organized as follows: Section 2 includes the formal description of the bargaining game and the definition of the equilibrium concept. In Section 3, we give necessary and sufficient conditions for a profile of stationary strategies to be an equilibrium in accordance with that concept. We also show that such an equilibrium exists. In Section 4, the main result is established: our non-cooperative support for the ANBS. Section 5 concludes.

2 The Bargaining Game

We consider the bargaining game $\Gamma(N, M, V)$. The set of players is denoted by N , and its members are indexed from 1 until n . The game is played for potentially infinitely many periods $t = 0, 1, 2, \dots$. In each period, one player acts as the proposer. In period $t = 0$, the proposer is determined in an arbitrary way. In all later periods, the proposer is chosen by a Markov chain. The probability distribution on the players in period $t > 0$ depends on the identity of the proposer in period $t - 1$. The entry m_{ji} of the matrix M is the probability that player j will propose in period t given that player i has proposed in period $t - 1$. All entries of M are nonnegative and for each $i \in N$, it is true that $\sum_{j=1}^n m_{ji} = 1$. The set V corresponds to all feasible payoffs. We denote $V \cap \mathbb{R}_+^n$ by V_+ . Our assumptions are as follows.

- (A1) The set V is closed, convex, and comprehensive from below. The origin lies in the interior of V . The set V_+ is bounded and all weakly Pareto-efficient points in V_+ are also strongly Pareto-efficient.
- (A2) The matrix M is irreducible.

We denote the interior and boundary of a set X by $\text{int}(X)$ and ∂X respectively. A vector η with $\|\eta\| = 1$ is said to be normal to the convex set V at a point $\bar{v} \in V$ if $(v - \bar{v})^\top \eta \leq 0$ for every $v \in V$. The set of all vectors η normal to V at \bar{v} is called the normal to V at \bar{v} .

(A3) There is a unique vector in the normal to V at every $v \in \partial V \cap \mathbb{R}_+^n$.

The assumption that all weakly Pareto-efficient points in V_+ are also strongly Pareto-efficient is essential to our results. As we show later in Section 3, this assumption implies that a proposal of a player i gives all other players their respective reservation payoffs. Thus a proposer always extracts the full surplus from all other players. Our assumptions with respect to V are similar to those in Merlo and Wilson (1995) and Miyakawa (2006).

The game proceeds as follows. In any period t , first the proposer is chosen in the aforementioned way. Next, the proposer offers a vector $v \in V$. Then, all players (including the proposer) decide sequentially whether to accept or reject the offer v , where for the sake of simplicity we assume that player i responds before player $i + 1$. We define the set $S(i)$ consisting of player i and all its successors by $S(i) = \{j \in N \mid j \geq i\}$. If all players have accepted the vector v in period t , the game ends and each player i receives a payoff of v_i . As soon as one player rejects v , period $t + 1$ starts with probability δ , and the game ends with probability $1 - \delta$. In the latter case, as well as in the case with perpetual disagreement, all players receive zero payoff. We assume that players maximize expected payoffs.

We denote by H_i^P the set of histories after which player i has to make a proposal and by H_i^R the set of histories after which player i has to respond to a proposal. Then, a strategy for player i is a map $s_i : H_i^P \cup H_i^R \rightarrow V \cup \{\text{Yes, No}\}$.

Player i 's strategy is stationary if the same proposal is made at all histories H_i^P and if the action taken at any history H_i^R depends only on the current proposal and the current proposer.¹

A Nash equilibrium is a profile of strategies from which no player has an incentive to unilaterally deviate. A subgame perfect Nash equilibrium (SPE) is a profile of strategies such that its restriction to any subgame is a Nash equilibrium of that subgame.

A stationary subgame perfect Nash equilibrium (SSPE) is a profile of stationary strategies which is an SPE.

3 Analysis of Stationary Equilibrium

In this section, we characterize the set of subgame perfect equilibria in stationary strategies and show that such equilibria exist. Theorem 3.8 gives the necessary and sufficient

¹This notion of stationarity is weaker than the notion of subgame consistency due to Harsanyi and Selten (1988), which implies that a player chooses the same action at any two nodes for which the continuation game is the same. For instance, suppose that rows i and j of M are identical. Then, the continuation games after rejection of player i 's proposal and that after rejection of player j 's proposal are identical. Yet, our definition allows A_k^i to be different from A_k^j for one or more players $k \in N$.

conditions for a strategy profile to be an SSPE and Theorem 3.13 shows that an SSPE exists.

The analysis in this section resembles Kalandrakis (2004), but some important differences should be noted: We admit only pure strategies and require unanimous agreement. Furthermore, we conclude rather than assume that agreement is immediate in SSPE and we do not impose assumptions on the behavior of players who are indifferent between acceptance and rejection of some proposal.

Consider a profile of stationary strategies. It can be described by an $n \times n$ -matrix Θ , where the entry θ_j^i is the payoff proposed to player j by player i , and a collection \mathcal{A} of n^2 acceptance sets, where the acceptance set A_j^i is the set of vectors in V which player j will accept when proposed by player i . The set of vectors in V proposed by player i and accepted by player j and his successors is $A_{S(j)}^i = \bigcap_{k \in S(j)} A_k^i$. We refer to $A^i = A_{S(1)}^i = \bigcap_{j \in N} A_j^i$ as the social acceptance set for proposer i .

Suppose that in period t , the proposal of player i is rejected. With probability $1 - \delta$ the game ends and all players receive zero payoff, and with probability δ period $t + 1$ is reached and play proceeds according to the profile (Θ, \mathcal{A}) of stationary strategies. The expected payoff to player j after rejection is $r_j^i(\Theta, \mathcal{A})$. Omitting the argument (Θ, \mathcal{A}) from the notation wherever possible, we refer to r_j^i as the reservation payoff of player j when i proposes.

Proposition 3.1 *The reservation payoff r^i belongs to $\text{int}(V)$.*

Proof: Conditional on the next period being reached, the payoffs are determined by a probability distribution on V (notice that also $0 \in V$), so expected payoffs belong to V since V is convex. Since with probability $1 - \delta$ the next period is not reached, these expected payoffs equal $\delta^{-1}r^i$, so $\delta^{-1}r^i \in V$. Since $0 \in \text{int}(V)$, the convex combination $(1 - \delta)0 + \delta\delta^{-1}r^i = r^i$ belongs to $\text{int}(V)$. \square

One implication of Proposition 3.1 is that a proposer always has the option to make a proposal that strictly exceeds the reservation payoff of every player.

Proposition 3.2 *In SSPE, for $j \in N$, if $v \in A_{S(j)}^i$, then $v_k \geq r_k^i$ for all $k \in S(j)$.*

Proof: Suppose that (Θ, \mathcal{A}) is a profile of stationary strategies such that $v \in A_{S(j)}^i$ but $v_k < r_k^i$ for some player $k \in S(j)$. Consider a history in H_k^i , where player k responds to the proposal v made by player i . At that history player k could deviate from (Θ, \mathcal{A}) by rejecting v . In that case, an expected payoff of r_k^i would result. Hence, this deviation is profitable and (Θ, \mathcal{A}) cannot be an SSPE. \square

Proposition 3.2 implies that for a vector of payoffs v to belong to the social acceptance set, it should satisfy $v_j \geq r_j^i$ for all $j \in N$.

Proposition 3.3 *In SSPE, for $j \in N$, if $v \in V$ satisfies $v_k > r_k^i$ for all $k \in S(j)$, then $v \in A_{S(j)}^i$.*

Proof: Suppose that (Θ, \mathcal{A}) is a profile of stationary strategies such that $v_k > r_k^i$ for all $k \in S(j)$ and $v \notin A_{S(j)}^i$. Suppose first that $v \notin A_n^i$. Consider a history in H_n^r , where player n responds to the proposal v made by player i . Then, player n could deviate from (Θ, \mathcal{A}) by accepting v . This would yield a payoff of $v_n > r_n^i$, a contradiction. Consequently, $v \in A_n^i$. We repeat the argument for players $n - 1, n - 2, \dots, j$ to establish the proposition. \square

Proposition 3.3 established a kind of converse of Proposition 3.2. One implication of this proposition is that a vector $v \in V$ that satisfies $v_j > r_j^i$ for all $j \in N$ belongs to the social acceptance set A_j^i .

Proposition 3.4 *In SSPE, each player's proposal θ^i lies in the social acceptance set A^i for proposer i .*

Proof: Suppose by way of contradiction that under some SSPE there is a player $i \in N$ such that $\theta^i \notin A^i$. Consider the subgame starting at a history where player i is the proposer. Since θ^i is rejected, r^i is the vector of expected payoffs by definition. By Proposition 3.1, $r^i \in \text{int}(V)$. Consequently, there exists $v \in V$ such that $v_j > r_j^i$ for all $j \in N$. By the previous proposition, $v \in A^i$. Hence, it would be a profitable deviation for player i to propose v instead of θ^i . \square

Proposition 3.5 *In SSPE, $\theta_j^i \geq 0$ and $r_j^i \geq 0$ for all $(i, j) \in N \times N$.*

Proof: Suppose by way of contradiction that (Θ, \mathcal{A}) is an SSPE and that $\theta_j^i < 0$ for some $(i, j) \in N \times N$. Consider a history where player j has to respond to the proposal θ^i . By Proposition 3.4, $\theta^i \in A^i$, so player j will receive a strictly negative payoff if play proceeds according to (Θ, \mathcal{A}) . But then, it would be a profitable deviation for player j to reject the proposal. Consequently, it holds that $\theta_j^i \geq 0$ for all $(i, j) \in N \times N$. It then follows that $r_j^i \geq 0$ for all $(i, j) \in N \times N$. \square

The next proposition establishes that an equilibrium proposal of any player gives all other players their respective reservation payoffs. Thus a proposer always extracts the entire surplus from the other players.

Proposition 3.6 *In SSPE, $\theta_j^i = r_j^i$ for all $(i, j) \in N \times N$ such that $i \neq j$.*

Proof: Since $\theta^i \in A^i$ by Proposition 3.4, Proposition 3.2 implies that $\theta_j^i \geq r_j^i$ for all $j \in N$. Suppose $\theta_k^i > r_k^i$ for some $k \in N$ such that $k \neq i$. Define the vector v as follows,

$$v_j = \begin{cases} \theta_j^i & \text{if } j = i \\ r_j^i & \text{if } j \neq i. \end{cases}$$

The vector v is clearly non-negative and it is in V , because $v \leq \theta^i$ and V is comprehensive. Furthermore, $v_k = r_k^i < \theta_k^i$, so the vector v is dominated by θ^i and is therefore not strongly Pareto-efficient. Since we assume that all weakly Pareto-efficient vectors of V_+ are also strongly Pareto-efficient, the vector v is not weakly Pareto-efficient. Thus, there exists a vector v' such that $v'_j > v_j$ for all $j \in N$.

We show now that $v' \in A^i$. Indeed, $v'_j > v_j = r_j^i$ for all $j \neq i$. And for player i we have the inequality $v'_i > v_i = \theta_i^i \geq r_i^i$. Thus we conclude that $v'_j > r_j^i$ for all $j \in N$. Proposition 3.3 now implies that $v' \in A^i$, as desired.

But then player i has a profitable deviation at any history where he is entitled to make a proposal, namely propose the vector v' rather than θ^i . Indeed, the vector v' is accepted and results in a payoff of $v'_i > \theta_i^i$ to player i . \square

Proposition 3.7 claims that all players make a proposal belonging to the boundary of V .

Proposition 3.7 *In SSPE, $\theta^i \in \partial V$ for each $i \in N$.*

Proof: Suppose by way of contradiction that there is a player $i \in N$ such that $\theta^i \in \text{int}(V)$. Equivalently, there exists $v \in V$ such that $v_j > \theta_j^i$ for all $j \in N$. By the immediate agreement property, $\theta^i \in A^i$. This implies that $v \in A^i$ as well. But then it would be a profitable deviation for player i to propose v rather than θ^i . \square

The previous propositions are collected in the following theorem.

Theorem 3.8 *If (Θ, \mathcal{A}) is an SSPE profile, then for all $i \in N$*

$$A_{S(j)}^i \subset \{v \in V | v_{j'} \geq \delta \sum_{k=1}^n m_{ki} \theta_j^k, \forall j' \in S(j)\}, \quad j \in N, \quad (1)$$

$$A_{S(j)}^i \supset \{v \in V | v_{j'} > \delta \sum_{k=1}^n m_{ki} \theta_j^k, \forall j' \in S(j)\} \cup \{\theta^i\}, \quad j \in N, \quad (2)$$

$$\theta^i \in \partial V \cap \mathbb{R}_+^n, \quad (3)$$

$$\theta_j^i = r_j^i = \delta \sum_{k=1}^n m_{ki} \theta_j^k, \quad j \in N \setminus \{i\}. \quad (4)$$

In what follows, we establish the converse; the four conditions of Theorem 3.8 characterize the set of SSPE. We first apply the well-known one-shot deviation principle to the game at hand. That is, we show that if there is a subgame where a player has some profitable deviation from a stationary strategy profile, then there must also be a subgame where this player has a profitable one-shot deviation. Here, a one-shot deviation in a subgame is a single deviation by the player at the root of the subgame. Next we show that no player has a profitable one-shot deviation.

Proposition 3.9 *Let (Θ, \mathcal{A}) be a stationary strategy profile satisfying (1)–(4). If there is a subgame such that some player has a profitable deviation, then there is a subgame where he has a profitable one-shot deviation.*

Proof: Consider a subgame starting in period t' and suppose that some, not necessarily stationary, strategy s^i leads to a higher payoff for player i than (θ^i, A^i) . Denote the gain to player i from the deviation by $\varepsilon > 0$. For $t > t'$, let $s^i(t)$ be the strategy that coincides with s^i in periods $t', \dots, t-1$ and with (θ^i, A^i) in periods $t, t+1, \dots$. In any round $t > t'$, we have by (1) that only proposals in V_+ can be accepted. Let \bar{v}_i be an upper bound on players i 's payoff when $v \in V_+$. Then the gain to player i from a deviation to $s_i(t)$ is at least $\varepsilon - \delta^{t-t'} \bar{v}^i / (1 - \delta)$. By choosing $t > t' + (\ln(\varepsilon) + \ln(1 - \delta) - \ln(\bar{v}^i)) / \ln(\delta)$ we have a profitable deviation $s^i(t)$ from (θ^i, A^i) that only involves a finite number of deviations.

If $s^i(t)$ involves a deviation by player i at a history $h \in H_i^r$ in period $t-1$, and $s^i(t)$ is a profitable deviation from (θ^i, A^i) in subgame $\Gamma(h)$, then the proposition holds. If $s^i(t)$ does not involve a profitable deviation by player i at any history $h \in H_i^r$ in period $t-1$, then the strategy $\bar{s}^i(t)$ that coincides with s^i in periods $t', \dots, t-2$ and at histories $h \in H_i^p$ in period $t-1$ and with (θ^i, A^i) otherwise, is a profitable deviation from (θ^i, A^i) in the subgame starting in period t' under consideration.

By a similar argument as in the previous paragraph, we can show that $\bar{s}^i(t)$ involves a profitable deviation by player i in subgame $\Gamma(h)$ with h a history in period $t-1$ belonging to H_i^p , or the strategy $s^i(t-1)$ is a profitable deviation from (θ^i, A^i) in the subgame starting in period t' under consideration. Repeating this argument a finite number of times leads to a subgame where player i has a profitable one-shot deviation. \square

Proposition 3.10 *Let (Θ, \mathcal{A}) be a stationary strategy profile satisfying (1)–(4). There is no subgame where a player has a profitable one-shot deviation.*

Proof: Consider the subgame at a history $h \in H_i^p$. Suppose player i has a one-shot deviation involving a proposal v^i different from θ^i . If v^i does not belong to A^i , it leads to a payoff r_i^i for player i . Since $\theta_j^i = r_j^i$ for all $j \neq i$ and $r^i \in \text{int}(V)$, the Pareto-efficiency of

θ^i implies $\theta_i^i > r_i^i$, so the deviation is not profitable. If v^i belongs to A^i , then, for $j \neq i$, $v_j^i \geq r_j^i = \theta_j^i$ by (1) and (4). Now $v_i^i \leq \theta_i^i$, since otherwise the Pareto-efficiency of θ^i would be violated. The deviation is not profitable.

Consider the subgame at a history $h \in H_i^i$. Let v denote the proposal being made, say by player j . Consider a one-shot deviation by player i . If $v \in A_i^j$ and $v \in A_{S(i)}^j$, then the deviation leads to a payoff r_i^j for player i , which is less than or equal to v_i , the payoff resulting from not deviating. If $v \in A_i^j$ and $v \notin A_{S(i)}^j$, then the deviation does not lead to a change in payoffs. If $v \notin A_i^j$ and $v \in \cap_{k \in S(i) \setminus \{i\}} A_k^j$, then the deviation leads to payoff v_i , which is less than or equal to r_i^j , the payoff resulting from not deviating. If $v \notin A_i^j$ and $v \notin \cap_{k \in S(i) \setminus \{i\}} A_k^j$, then the deviation does not lead to a change in payoff. \square

The previous two propositions imply that the strategies conforming to the conditions of Theorem 3.8 are subgame-perfect. Since they are also stationary, we have the following

Theorem 3.11 *The strategy profile (Θ, \mathcal{A}) fulfills the four conditions of Theorem 3.8 if and only if it is an SSPE.*

If a matrix Θ of proposals is part of an SSPE, then it is part of many SSPE's. This inessential multiplicity has two sources. First, if a responding player is proposed exactly the reservation payoff, then our characterization restricts behavior only if the proposal on the table is the equilibrium proposal. This is reflected in the fact that each point on the boundary of a social acceptance set (except the relevant equilibrium proposal) may or may not be an element of that set in SSPE. Second, if a proposal lies outside a social acceptance set, it is indeterminate which player will reject the proposal. Consider for example the case where $N = \{1, 2, 3, 4\}$ and $r^1 = (1, 1, 1, 1)$. Now suppose that player 1 has proposed $v = (2, 2, 0, 0)$ in some subgame. Since $v_3 < r_3^1$ and $v_4 < r_4^1$, Proposition 3.2 implies that $v \notin A^1$, and by Proposition 3.4 $v \neq \theta^1$. It is also true that $v \notin A_4^1$: If the node where player 4 has to respond is reached, that player effectively chooses between a payoff of 1 and a payoff of 0, so SSPE requires rejection of the proposal. However, the SSPE characterization leaves indeterminate whether players 1, 2, and 3 will accept or reject v . Consequently, there is an SSPE for any configuration of responses by these players. In particular, player 3 may accept v in SSPE although $v_3 < r_3^1$, and player 2 may reject v although $v_2 > r_2^1$. This reasoning even extends to player 1: In SSPE, it is possible that $v \notin A_1^1$ although $v_1 > r_1^1$ and player 1 is the proposer. However, this does not mean that player 1 may reject his own proposal on the equilibrium path, since the SSPE characterization requires the specific proposal θ^1 to be made and immediately accepted by all players.

Proposition 3.12 *Given $v \in V_+$, let $W(v) = \{x \in V | x \geq v\}$. For each $i \in N$ there exists a unique $v' \in \partial V \cap \mathbb{R}_+^n$ such that $v'_j = v_j$ for all $j \in N \setminus \{i\}$. Furthermore,*

$$\{v'\} = \arg \max_{x \in W(v)} x_i.$$

Proof: Take any

$$v' \in \arg \max_{x \in W(v)} x_i.$$

We show that $v' \in \partial V \cap \mathbb{R}_+^n$ and $v'_j = v_j$ for all $j \in N \setminus \{i\}$. Indeed, $v' \in \mathbb{R}_+^n$, because $W(v) \subset \mathbb{R}_+^n$. Furthermore, v' is a boundary point of V , for if it was an interior point, there would exist a vector $\hat{v} \in V$ such that $\hat{v}_i > v'_i$ for all $i \in N$, contradicting the choice of v' as a maximizer of x_i over all points $x \in W(v)$.

Suppose $v'_k > v_k$ for some $k \in N \setminus \{i\}$. Define the vector x by the equation

$$x_j = \begin{cases} v'_j & \text{if } j \neq k \\ v_k & \text{if } j = k. \end{cases}$$

Then $v' \geq x$, so x is an element of V , because V is comprehensive from below. Furthermore, $x \geq v \geq 0$. Since $v' \geq x$ with the strict inequality for component k , x is not strongly Pareto-efficient. Since we assume that all weakly Pareto-efficient vectors of V_+ are also strongly Pareto-efficient, x is not weakly Pareto-efficient, so there exists $\hat{v} \in V$ such that $\hat{v}_j > x_j$ for all $j \in N$. But then $\hat{v} \geq x \geq v$, so that $\hat{v} \in W(v)$, and $\hat{v}_i > x_i = v'_i$, contradicting the choice of v' as a maximizer of x_i over all points x in $W(v)$.

Suppose now that there exists a vector $v'' \in \partial V \cap \mathbb{R}_+^n$ other than v' such that $v''_j = v_j$ for all $j \in N \setminus \{i\}$. Without loss of generality assume $v''_i > v'_i$. Then the vector v'' is not strongly Pareto-efficient. Since we assume that each weakly Pareto-efficient vector is also strongly Pareto-efficient, v'' is not weakly Pareto-efficient. But then v'' is in the interior of V , a contradiction. \square

Theorem 3.13 *An SSPE exists.*

Proof: We identify the set of SSPE of Γ with the set of fixed points of a map and use Brouwer's fixed point theorem to establish existence. Consider a non-negative matrix Θ with each column θ^i of Θ in the set V_+ . Let $\chi^i(\Theta)$ be a vector satisfying the following conditions:

$$\chi^i(\Theta) \in \partial V \cap \mathbb{R}_+^n \text{ and } \chi_j^i(\Theta) = \delta \sum_{k=1}^n m_{ki} \theta_j^k \text{ for all } j \neq i.$$

By Proposition 3.12, $\chi^i(\Theta)$ is uniquely determined. Define now the function $\chi : V_+ \rightarrow V_+$ by letting $\chi(\Theta) = (\chi^1(\Theta), \dots, \chi^n(\Theta))$. By the characterization of SSPE, if $\chi(\Theta) = \Theta$ for a particular $\Theta \in V_+$, then there is at least one SSPE with Θ as the matrix of proposed payoffs. The set $V_+ = V \cap \mathbb{R}_+^n$ contains the origin, it is convex and closed, as it is the intersection of two convex and closed sets, and it is bounded. Since the maximum theorem can be used to show that χ is continuous, Brouwer's fixed point theorem implies that it has a fixed point. \square

4 The Limit Equilibrium

Our proofs so far did not rely on Assumptions A2 and A3. They will be needed for the results of this section. Since the matrix M is irreducible, it has a unique stationary distribution denoted by μ . Recall that the stationary distribution μ is a probability distribution on the set of players N satisfying the equation $M\mu = \mu$. Furthermore, irreducibility of M implies that all states occur with positive probability under the stationary distribution, that is $\mu_i > 0$ for each $i \in N$. If the matrix M was reducible, the state space of M could be partitioned into several communicating classes. In this case, one obtains results analogous to those in the sequel within each communicating class.

Theorem 4.2 below is the main result of the paper. As the continuation probability goes to one, along any sequence of stationary subgame perfect equilibria of Γ , the equilibrium proposal of all players converges to the same limit. This common limit is the asymmetric Nash bargaining solution weighted by the stationary distribution μ , denoted μ -ANBS.

Definition 4.1 The *asymmetric Nash product* with weights μ is the function $\rho : V \rightarrow \mathbb{R}$ defined by

$$\rho(v) = \prod_{i \in N} (v_i)^{\mu_i}.$$

The μ -ANBS is the unique maximizer of the function ρ on the set V .

Theorem 4.2 Let $\{\delta_m\}$ be a sequence of continuation probabilities in $[0, 1)$ converging to 1. For each m , let $\Theta(\delta_m)$ be a matrix of proposals in some SSPE of the game Γ with continuation probability δ_m . Then the limits $\lim \theta^i(\delta_m)$ exist for each $i \in N$. All limits are equal to the μ -ANBS.

Let δ_m and $\Theta(\delta_m)$ be as in Theorem 4.2. The sequence $\{\Theta(\delta_m)\}$ has a convergent subsequence, as it lies in the compact set V_+ . For the remainder of this section, we will fix any such convergent subsequence and denote its limit by $\bar{\Theta}$. Since the convergent

subsequence considered is arbitrary, to prove Theorem 4.2 it is sufficient to show that each column of the matrix $\bar{\Theta}$ is the μ -ANBS.

We now give a brief overview of the argument. First we show that along the sequence $\{\Theta(\delta_m)\}$ of equilibria the proposals of all players converge to a common limit, say the point $\bar{\theta} \in V$. We then compute the tangent space to the set ∂V at the point $\bar{\theta}$ by considering the pairwise differences of the equilibrium proposals of players i and n , and show they converge to zero at the same speed as $1 - \delta_m$. In fact, we are able to compute explicitly the limits of the vectors $(\theta^i(\delta_m) - \theta^n(\delta_m))/(1 - \delta_m)$, which are then shown to span the tangent space to the set ∂V at the point $\bar{\theta}$. Using this result we show next that the tangent space at $\bar{\theta}$ is orthogonal to the gradient of the asymmetric Nash product with weights μ , thereby showing that $\bar{\theta}$ is the μ -ANBS.

Proposition 4.3 *All columns of $\bar{\Theta}$ are identical.*

Proof: For any $i \neq j$, it follows from the SSPE characterization that

$$\bar{\theta}_j^i = \sum_{k=1}^n m_{ki} \bar{\theta}_j^k.$$

Suppose that, contrary to the proposition, not all limit proposals $\bar{\theta}^1, \dots, \bar{\theta}^n$ are the same. Let $j \in N$ be such that in the limit not all players propose the same to player j , and choose $\bar{\theta}_j^\diamond$ to be either $\min_{i \in N} \{\bar{\theta}_j^i\}$ or $\max_{i \in N} \{\bar{\theta}_j^i\}$, whichever is *not* equal to $\bar{\theta}_j^j$. Define $N^\diamond = \{i \in N \mid \bar{\theta}_j^i = \bar{\theta}_j^\diamond\}$. For any $i \in N^\diamond$, we have

$$\bar{\theta}_j^\diamond = \bar{\theta}_j^\diamond \sum_{k \in N^\diamond} m_{ki} + \sum_{k \in N \setminus N^\diamond} m_{ki} \bar{\theta}_j^k,$$

which is equivalent to

$$\bar{\theta}_j^\diamond \sum_{k \in N \setminus N^\diamond} m_{ki} = \sum_{k \in N \setminus N^\diamond} m_{ki} \bar{\theta}_j^k.$$

Suppose first that for some $i \in N^\diamond$,

$$\sum_{k \in N \setminus N^\diamond} m_{ki} > 0.$$

Then

$$\bar{\theta}_j^\diamond = \frac{\sum_{k \in N \setminus N^\diamond} m_{ki} \bar{\theta}_j^k}{\sum_{k \in N \setminus N^\diamond} m_{ki}},$$

which contradicts the fact that either $\bar{\theta}_j^\diamond < \bar{\theta}_j^k$ for all $k \in N \setminus N^\diamond$ or $\bar{\theta}_j^\diamond > \bar{\theta}_j^k$ for all $k \in N \setminus N^\diamond$. Therefore, $m_{ki} = 0$ for all $(i, k) \in N^\diamond \times (N \setminus N^\diamond)$ and thus N^\diamond is an absorbing set. Since $N^\diamond \neq N$, this contradicts the irreducibility assumption on M . \square

We denote a column of $\bar{\Theta}$ by $\bar{\theta}$. For $i \in N \setminus \{n\}$, we define

$$d^i(\delta_m) = \frac{\theta^i(\delta_m) - \theta^n(\delta_m)}{1 - \delta_m}.$$

Let $D(\delta_m)$ be the $n \times (n-1)$ -matrix with columns $d^1(\delta_m), \dots, d^{n-1}(\delta_m)$.

The rest of the proof is organized as follows. In Proposition 4.5 we compute the limits of $d^i(\delta_m)$ as m goes to infinity. Proposition 4.6 establishes that the limit of $\{d^i(\delta_m)\}$ belongs to the tangent space to ∂V at $\bar{\theta}$. We then proceed to show in Proposition 4.8 that the limits of $d^1(\delta_m), \dots, d^{n-1}(\delta_m)$ are linearly independent and thus span the tangent space to ∂V at $\bar{\theta}$. And finally, Proposition 4.9 establishes that the gradient of the Nash product with weights μ at the point $\bar{\theta}$ is orthogonal to the tangent space of V , thereby showing that $\bar{\theta}$ is the μ -ANBS.

For a matrix A , we denote by A_{-i} and A^{-i} the matrix A without its i^{th} row and column, respectively. We write $\mathbf{1}$ for a column vector of ones and I for the identity matrix.

For $j \in N$, we define the matrix $L(j)$ by

$$L(j) = [M - I]_{-n}^{-j}.$$

Thus $L(j)$ is the $(n-1) \times (n-1)$ -matrix obtained from $M - I$ by deleting column j and row n . Proposition 4.4 is an auxiliary result used in the proof of Proposition 4.5.

Proposition 4.4 *The matrix $L(j)$ is invertible for all $j \in N$.*

Proof: Suppose $L(j)$ is singular. Let a be a non-zero vector such that $[M - I]_{-n}^{-j}a = 0$. Since the elements in any column of the matrix $M - I$ add up to zero, we also have the equation $(M - I)_n^{-j}a = 0$, so $[M - I]^{-j}a = 0$. By using $[M - I]^{-j} = (M - I)I^{-j}$ and defining $b = I^{-j}a$, we see that $(M - I)b = 0$. Thus the vector b is an eigenvector of M associated with eigenvalue 1. By the Perron–Frobenius theorem, any non-zero eigenvector of M associated with eigenvalue 1 is a strictly positive vector. However, since $b_j = 0$, we have obtained a contradiction. Consequently, the matrix $L(j)$ is invertible. \square

Proposition 4.5 *The sequence $\{D(\delta_m)\}$ of matrices converges to the matrix \bar{D} with rows given by $\bar{d}_j = \bar{\theta}_j \mathbf{1}^\top L^{-1}(j)$ for $j \in N$.*

Proof: We fix m and denote δ_m by δ , $\theta^i(\delta_m)$ by θ^i , and $d^i(\theta^m)$ by d^i .

For each $j \in N$ and $i \in N \setminus \{j, n\}$,

$$\begin{aligned} d_j^i(1 - \delta) &= \theta^i - \theta^n = \delta \sum_{k=1}^n m_{ki} \theta_j^k - \theta_j^n \\ &= \delta \sum_{k=1}^n m_{ki} (\theta_j^k - \theta_j^n) + \delta \theta_j^n - \theta_j^n, \end{aligned}$$

where we use (4) for the second equality, so

$$d_j^i = \delta \sum_{k=1}^{n-1} m_{ki} d_j^k - \theta_j^n.$$

We have found that

$$\theta_j^n = \delta \sum_{k \notin \{i, n\}} m_{ki} d_j^k + (\delta m_{ii} - 1) d_j^i, \quad j \in N, \quad i \in N \setminus \{j, n\}. \quad (5)$$

Similarly, for $j \neq n$,

$$\begin{aligned} d_j^j(1 - \delta) &= \theta_j^j - \delta \sum_{k=1}^n m_{kn} \theta_j^k \\ &= \theta_j^j - \delta \sum_{k=1}^n m_{kn} (\theta_j^k - \theta_j^n) - \delta \theta_j^n \\ &= \theta_j^j - \theta_j^n - \delta \sum_{k=1}^n m_{kn} (\theta_j^k - \theta_j^n) + (1 - \delta) \theta_j^n, \end{aligned}$$

where we use (4) for the second inequality, so

$$d_j^j = d_j^j - \delta \sum_{k=1}^{n-1} m_{kn} d_j^k + \theta_j^n.$$

We have found that

$$\theta_j^n = \delta \sum_{k=1}^{n-1} m_{kn} d_j^k, \quad j \in N \setminus \{n\}. \quad (6)$$

We write (5)–(6) in vector–matrix notation as

$$\theta_j^n \mathbf{1}^\top = d_j (\delta M - I)_{-n}^{-j}, \quad j \in N.$$

The matrix $(M - I)_{-n}^{-j}$ is invertible by Proposition 4.4, and so is the matrix $(\delta M - I)_{-n}^{-j}$ for δ close enough to one. Thus, for every $j \in N$, we can solve the above system for d_j as

$$d_j = \theta_j^n \mathbf{1}^\top [(\delta M - I)_{-n}^{-j}]^{-1}.$$

As δ_m goes to one, the sequence $\theta_j^n(\delta_m)$ converges to $\bar{\theta}_j$ by Proposition 4.3. Thus the sequence $d_j(\delta_m)$ converges to $\bar{\theta}_j \mathbf{1}^\top L^{-1}(j)$, as desired. \square

Proposition 4.5 expresses each row j of the matrix \bar{D} as the sum of the rows of the matrix $L^{-1}(j)$ multiplied by the scalar $\bar{\theta}_j$.

We show now that each column of the matrix \bar{D} is orthogonal to the normal vector of V at the point $\bar{\theta}$, which is unique by Assumption A3. This is equivalent to saying that each column of the matrix \bar{D} belongs to the tangent space of ∂V at $\bar{\theta}$. We let $\text{span}(\bar{D})$ denote the column span of the matrix \bar{D} .

Proposition 4.6 *It holds that $\text{span}(\bar{D})$ is orthogonal to the normal vector of V at $\bar{\theta}$.*

Proof: Let $\eta^i(\delta_m)$ denote the normal vector of V at the point $\theta^i(\delta_m)$. Since $\{\theta^i(\delta_m)\}$ converges to $\bar{\theta}$, the sequence $\{\eta^i(\delta_m)\}$ converges to $\bar{\eta}$, the normal vector to the set V at the point $\bar{\theta}$. By the definition of the normal vector,

$$\eta^n(\delta_m)^\top (\theta^i(\delta_m) - \theta^n(\delta_m)) \leq 0 \text{ and } \eta^i(\delta_m)^\top (\theta^i(\delta_m) - \theta^n(\delta_m)) \geq 0.$$

Dividing by $1 - \delta_m$ and passing to the limit yields the inequalities $\bar{\eta}^\top \bar{d}^i \leq 0$ and $\bar{\eta}^\top \bar{d}^i \geq 0$, therefore $\bar{\eta}^\top \bar{d}^i = 0$, as desired. \square

Propositions 4.7 and 4.8 address the dimension of $\text{span}(\bar{D})$. We show that the columns of \bar{D} are linearly independent, thus establishing that $\text{span}(\bar{D})$ equals the tangent space of ∂V at $\bar{\theta}$.

For $j \in N$, let K_j be the sum of the rows of the matrix $L^{-1}(j)$, thus

$$K_j = \mathbf{1}^\top L^{-1}(j).$$

Define K as the $n \times (n - 1)$ -matrix with rows K_j . Proposition 4.7 expresses all rows of K in terms of rows of $L^{-1}(n)$ and the stationary distribution μ induced by M .

Proposition 4.7 *Any combination of $n - 1$ distinct rows of the matrix K is linearly independent. Furthermore,*

$$K_j = \mathbf{1}^\top L^{-1}(j) = \mathbf{1}^\top L^{-1}(n) - \frac{1}{\mu_j} (L^{-1}(n))_j, \quad j \in N \setminus \{n\}.$$

Proof: We define $x = [M - I]_{-n}^n$. Consider some $j \in N \setminus \{n\}$. It can be verified by a direct computation that

$$L^{-1}(j) = \begin{bmatrix} (L^{-1}(n))_1 - \frac{(L^{-1}(n)x)_1}{(L^{-1}(n)x)_j} (L^{-1}(n))_j \\ \vdots \\ (L^{-1}(n))_{j-1} - \frac{(L^{-1}(n)x)_{j-1}}{(L^{-1}(n)x)_j} (L^{-1}(n))_j \\ (L^{-1}(n))_{j+1} - \frac{(L^{-1}(n)x)_{j+1}}{(L^{-1}(n)x)_j} (L^{-1}(n))_j \\ \vdots \\ (L^{-1}(n))_{n-1} - \frac{(L^{-1}(n)x)_{n-1}}{(L^{-1}(n)x)_j} (L^{-1}(n))_j \\ \frac{1}{(L^{-1}(n)x)_j} (L^{-1}(n))_j \end{bmatrix}.$$

The formula above is well-known in linear programming and is used to compute the simplex tableau following from a change in basis variables. By definition of the stationary distribution we have

$$L(n)\mu_{-n} + x\mu_n = 0.$$

We multiply this expression by $L^{-1}(n)$ and rearrange to obtain

$$L^{-1}(n)x = -\mu_{-n}/\mu_n.$$

By substitution, we find that

$$L^{-1}(j) = \begin{bmatrix} (L^{-1}(n))_1 - \frac{\mu_1}{\mu_j}(L^{-1}(n))_j \\ \vdots \\ (L^{-1}(n))_{j-1} - \frac{\mu_{j-1}}{\mu_j}(L^{-1}(n))_j \\ (L^{-1}(n))_{j+1} - \frac{\mu_{j+1}}{\mu_j}(L^{-1}(n))_j \\ \vdots \\ (L^{-1}(n))_{n-1} - \frac{\mu_{n-1}}{\mu_j}(L^{-1}(n))_j \\ -\frac{\mu_n}{\mu_j}(L^{-1}(n))_j \end{bmatrix}.$$

Summing up the rows of $L^{-1}(j)$ we get

$$\mathbf{1}^\top L^{-1}(j) = \sum_{i \in N \setminus \{j, n\}} (L^{-1}(n))_i + \frac{\mu_j - 1}{\mu_j}(L^{-1}(n))_j = \mathbf{1}^\top L^{-1}(n) - \frac{1}{\mu_j}(L^{-1}(n))_j.$$

Therefore,

$$K_{-n} = [\mathbf{1}\mathbf{1}^\top - C]L^{-1}(n),$$

where C is the $(n-1)$ -diagonal matrix with element $1/\mu_i$ in column i .

The matrix $[\mathbf{1}\mathbf{1}^\top - C]$ is non-singular. Suppose not, then there is $y \neq 0$ such that $[\mathbf{1}\mathbf{1}^\top - C]y = 0$. It follows that $\mathbf{1}\mathbf{1}^\top y = Cy = (y_1/\mu_1, \dots, y_{n-1}/\mu_{n-1})^\top$, from which it follows in particular that $\mathbf{1}^\top y \neq 0$. By pre-multiplying the last equality with the row vector $(\mu_1, \dots, \mu_{n-1})$, we find that $(1 - \mu_n)\mathbf{1}^\top y = \mathbf{1}^\top y$, a contradiction since $\mu_n > 0$. Consequently, the matrix $[\mathbf{1}\mathbf{1}^\top - C]$ is non-singular.

It follows that K_{-n} is non-singular. Since the labeling of players is arbitrary, we have shown that any combination of $n-1$ distinct rows of the matrix K is linearly independent. \square

Proposition 4.8 *It holds that $\bar{\theta}_i > 0$ for all $i \in N$. The column span of the matrix \bar{D} has dimension $n - 1$.*

Proof: We know that $\bar{\theta}_i \geq 0$ for each $i \in N$. Partition N into the set N_0 of players i such that $\bar{\theta}_i = 0$ and the set N_+ of players i such that $\bar{\theta}_i > 0$.

Suppose that the set N_0 is non-empty, so that the set N_+ consists of at most $n - 1$ elements. We show first that $\bar{\theta} \in \text{span}(\bar{D})$ by constructing a vector $z \in \mathbb{R}^{n-1}$ such that $\bar{D}z = \bar{\theta}$. Since the rows K_i of the matrix K corresponding to the elements i of the set N_+ are linearly independent by Proposition 4.7, there exists a vector $z \in \mathbb{R}^{n-1}$ such that $K_i z = 1$ for all $i \in N_+$. Then $\bar{d}_i z = \bar{\theta}_i K_i z = \bar{\theta}_i$ for all $i \in N_+$. Trivially, z also satisfies the equations $\bar{d}_i z = \bar{\theta}_i K_i z = 0 = \bar{\theta}_i$ for each $i \in N_0$.

Let $\bar{\eta}$ be the normal vector to V at the point $\bar{\theta}$. Since $\bar{\eta}$ is orthogonal to $\text{span}(\bar{D})$ by Proposition 4.6, we have $\bar{\eta}^\top \bar{\theta} = 0$. Since zero is in the interior of V by assumption, the vector $\varepsilon \bar{\eta}$ is in the set V for $\varepsilon > 0$ small enough. But then we have the inequality $\bar{\eta}^\top (\varepsilon \bar{\eta} - \bar{\theta}) = \varepsilon (\bar{\eta}^\top \bar{\eta}) > 0$, contradicting the definition of a normal vector. Consequently, the set N_0 is empty. We have established the first part of the proposition.

To prove the second part of the proposition, notice that \bar{D} can be written as the product TK , where T is a diagonal matrix with $\bar{\theta}_i$ in column i . Since $\bar{\theta}_i > 0$ for each $i \in N$, the matrix T has full rank n , and the matrix K has rank $n - 1$ by Proposition 4.7. This establishes the second part of the proposition. \square

We now establish that the gradient of the logarithm of the asymmetric Nash product with weights μ is orthogonal to the column span of the matrix \bar{D} . We observe that

$$\ln \rho(v) = \sum_{i \in N} \mu_i \ln(v_i).$$

Proposition 4.9 *It holds that $\text{span}(\bar{D})$ is orthogonal to the gradient of the function $\ln \rho$ at the point $\bar{\theta}$.*

Proof: The gradient of $\ln \rho$ at $\bar{\theta}$ is the vector g given by $g_j = \mu_j / \bar{\theta}_j$, $j \in N$. We have the following chain of equations

$$\begin{aligned} \sum_{j \in N} g_j \bar{d}_j &= \sum_{j \in N} \mu_j \mathbf{1}^\top L^{-1}(j) \\ &= \sum_{j \in N \setminus \{n\}} \mu_j [\mathbf{1}^\top L^{-1}(n) - \frac{1}{\mu_j} (L^{-1}(n))_j] + \mu_n \mathbf{1}^\top L^{-1}(n) \\ &= \mathbf{1}^\top L^{-1}(n) - \sum_{j \in N \setminus \{n\}} (L^{-1}(n))_j \\ &= \mathbf{1}^\top L^{-1}(n) - \mathbf{1}^\top L^{-1}(n) = 0, \end{aligned}$$

where the first equality uses the result of Proposition 4.5, and the second one Proposition 4.7. This establishes the proposition. \square

The proof of Theorem 4.2 is now immediate. The column span of the matrix \bar{D} is orthogonal to the normal vector of V at $\bar{\theta}$ by Proposition 4.6, and at the same time it is orthogonal to the gradient of $\ln \rho$ at $\bar{\theta}$ by Proposition 4.9. Since $\text{span}(\bar{D})$ has dimension $n-1$ (Proposition 4.8), it follows at once that the gradient of the function $\ln \rho$ is proportional to the normal vector to V at $\bar{\theta}$. Hence, the point $\bar{\theta}$ is the maximizer of the function $\ln \rho$ on the set V , as well as the maximizer of the function ρ on the set V .

5 Conclusion

In this paper we have provided further non-cooperative support to the asymmetric Nash bargaining solution. We demonstrate that existing results are instances of a much more generally valid principle. We consider a bargaining process involving any number of players, an arbitrary irreducible Markov process that determines the selection of the proposer, and any set of feasible payoffs that is bounded, convex, and has a smooth boundary. As long as no agreement is reached, negotiations break down with some fixed probability.

We fully characterize the set of subgame perfect equilibria in stationary strategies. We show that at least one such equilibrium exists and argue that in general there are many such equilibria. We continue by studying the limit of an arbitrary sequence of equilibria when the probability of breakdown goes to zero. We establish that in the limit all players make the same proposal. Moreover, this proposal is the same as the one corresponding to the asymmetric Nash bargaining solution, where the weights in the Nash product are equal to the stationary distribution of the Markov process that determines the selection of the proposer.

One implication is that if players are selected as proposer in some fixed order, then the symmetric Nash bargaining solution is achieved in the limit. This can be seen as a generalization of alternating offer bargaining to more than two players. Another implication is that if players are selected according to time-invariant probabilities, these probabilities are equal to the weights in the Nash product. The symmetric Nash bargaining solution would again result if the time-invariant probabilities are uniform.

It is noteworthy that the bargaining power of the players is only affected by the stationary distribution of the proposer selection process. The particular shape of the set of feasible payoffs is irrelevant for the weights of the players in the Nash product, as are the particular probabilities by which the proposer in the next period is chosen conditional on the current proposer.

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